

# Static Memory Leak Detection Using Full-Sparse Value-Flow Analysis



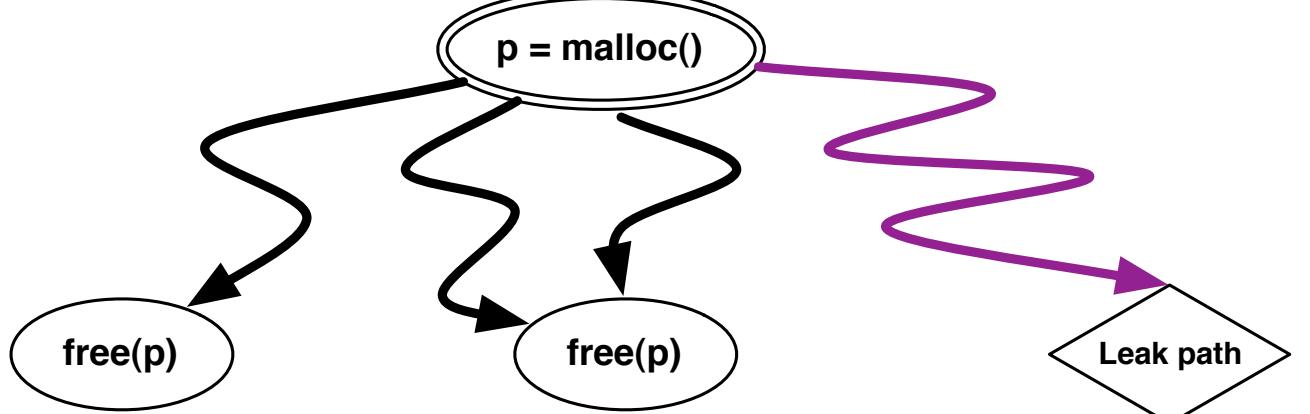
Yulei Sui, Ding Ye, and Jingling Xue

{YSUI,DYE,JINGLING}@CSE.UNSW.EDU.AU

SCHOOL OF COMPUTER SCIENCE AND ENGINEERING, UNSW

## PROBLEM

To find memory leaks **statically** in a program (without actually running it), a leak analysis reasons about a source-sink property: every object created at an allocation site (a source) must eventually reach a free site (a sink) during any execution of the program.



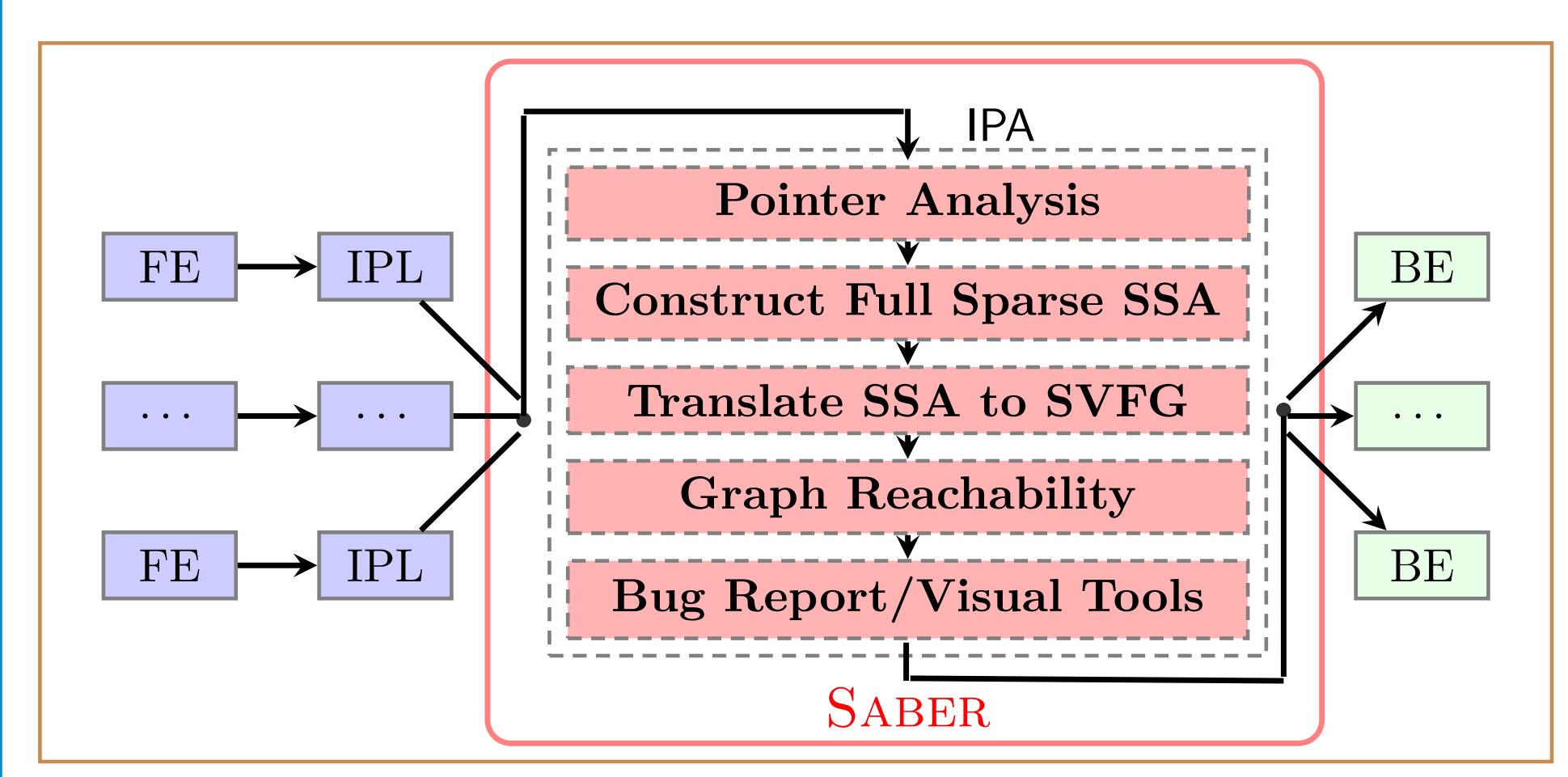
## CONTRIBUTIONS

1. The first to find memory leaks by using a full-sparse value-flow analysis to track the flow of values through all memory locations.
2. The first to Leverage recent advances on sparse pointer analysis in a major client application.
3. Effective at detecting 211 leaks at a false positive rate of 18.5% in the 15 SPEC2000 and 5 open-source C programs (totalling in 2324.1 KLOC).

| Leak Detector | Speed (LOC/sec) | Bug Count | False Positive Rate (%) |
|---------------|-----------------|-----------|-------------------------|
| CONTRADICTION | 300             | 26        | 56                      |
| CLANG         | 400             | 27        | 25                      |
| SPARROW       | 720             | 81        | 16                      |
| FASTCHECK     | 37,900          | 59        | 14                      |
| <b>SABER</b>  | <b>10,220</b>   | <b>83</b> | <b>19</b>               |

Comparing SABER with other detectors using SPEC2000

## OUR TOOL FRAMEWORK



We have implemented our tool **SABER** in Open64, an open-source industry-strength compiler, at IPA (interprocedural analysis module) containing four phases.

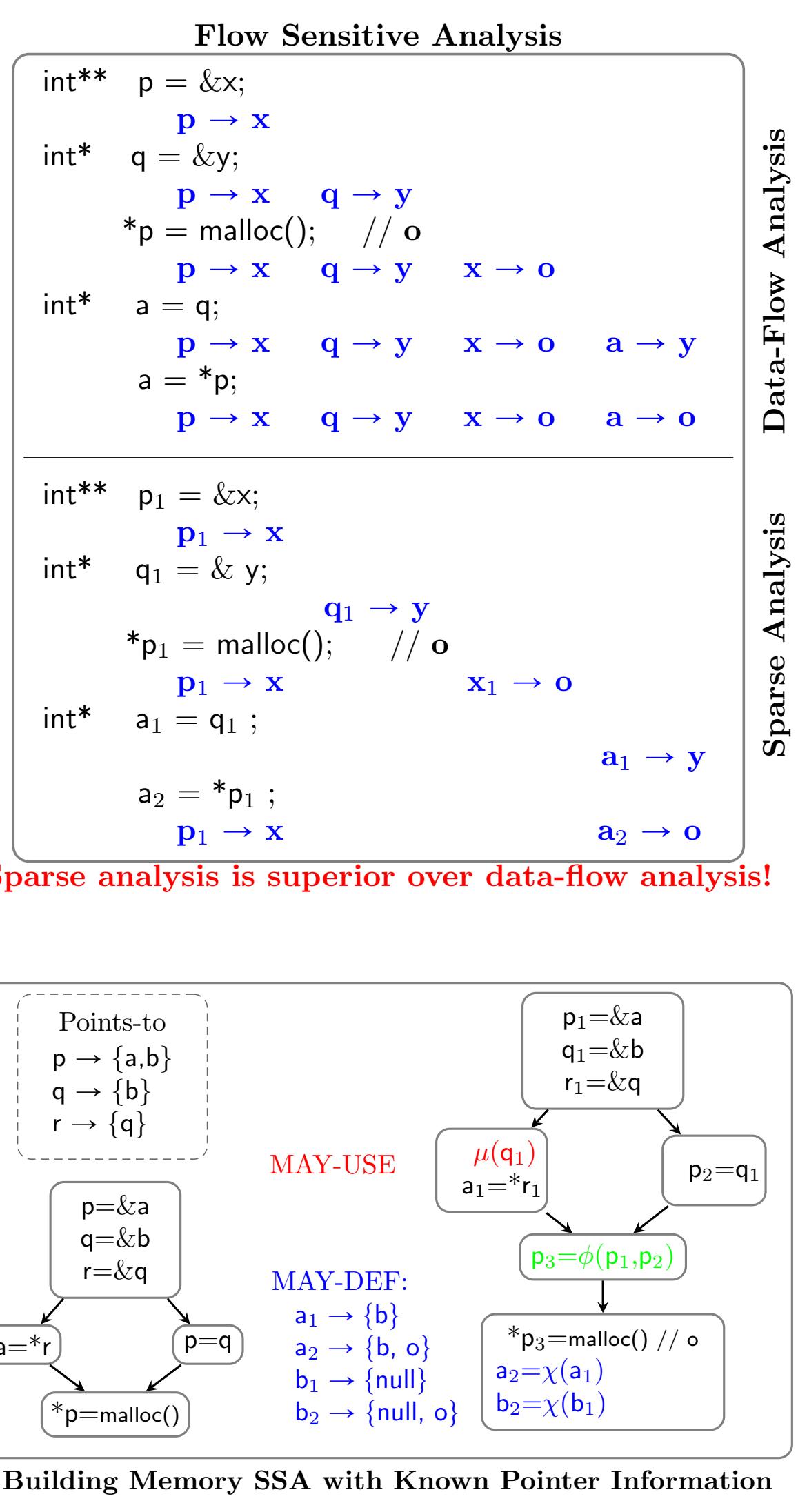
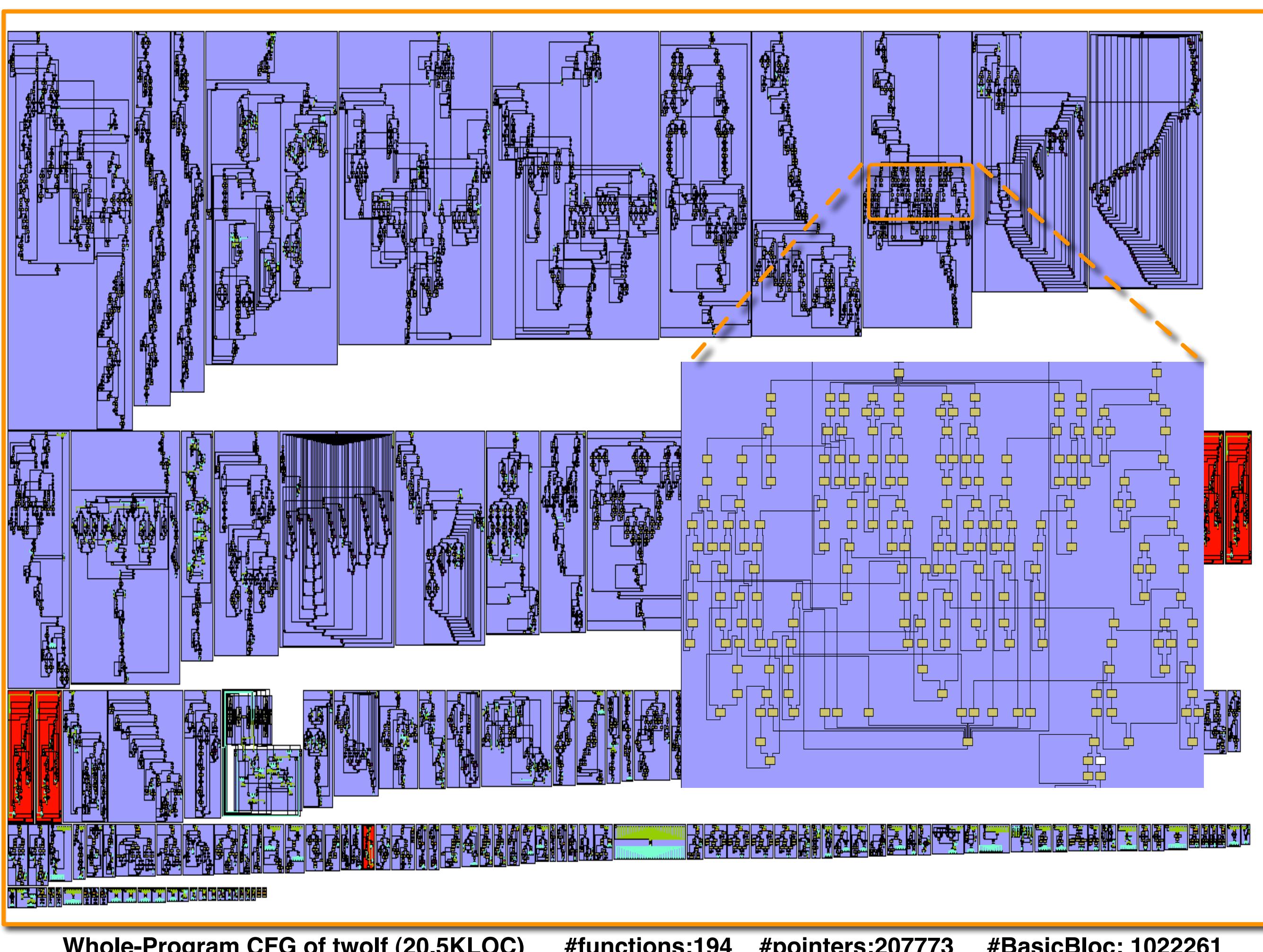
**IPA:** Global analysis by combining information from IPL

**IPL:** Summary information local to a function

**FE:** Compiler Front End

**BE:** Compiler Backend End

## BACKGROUND



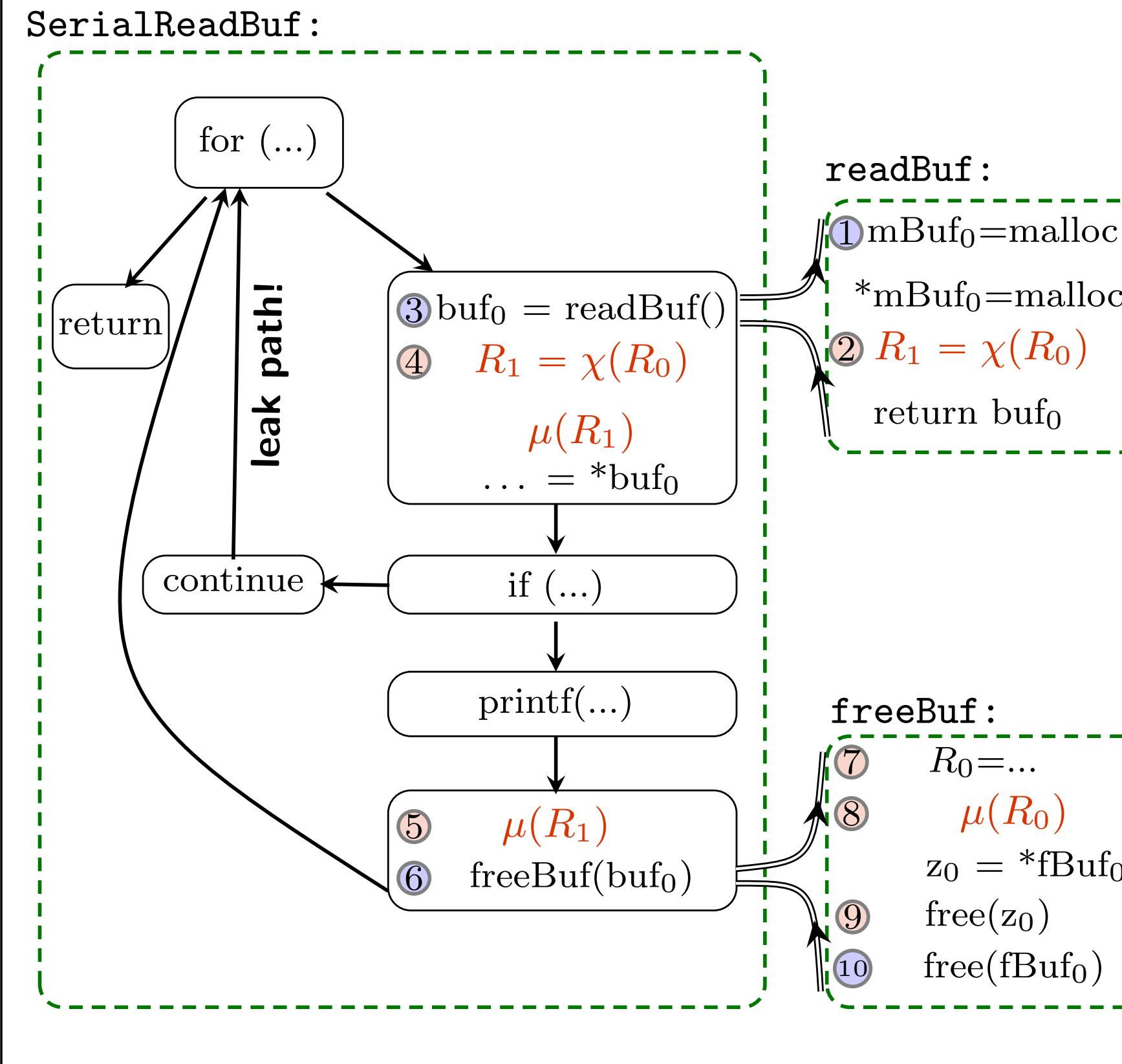
## AN EXAMPLE

```

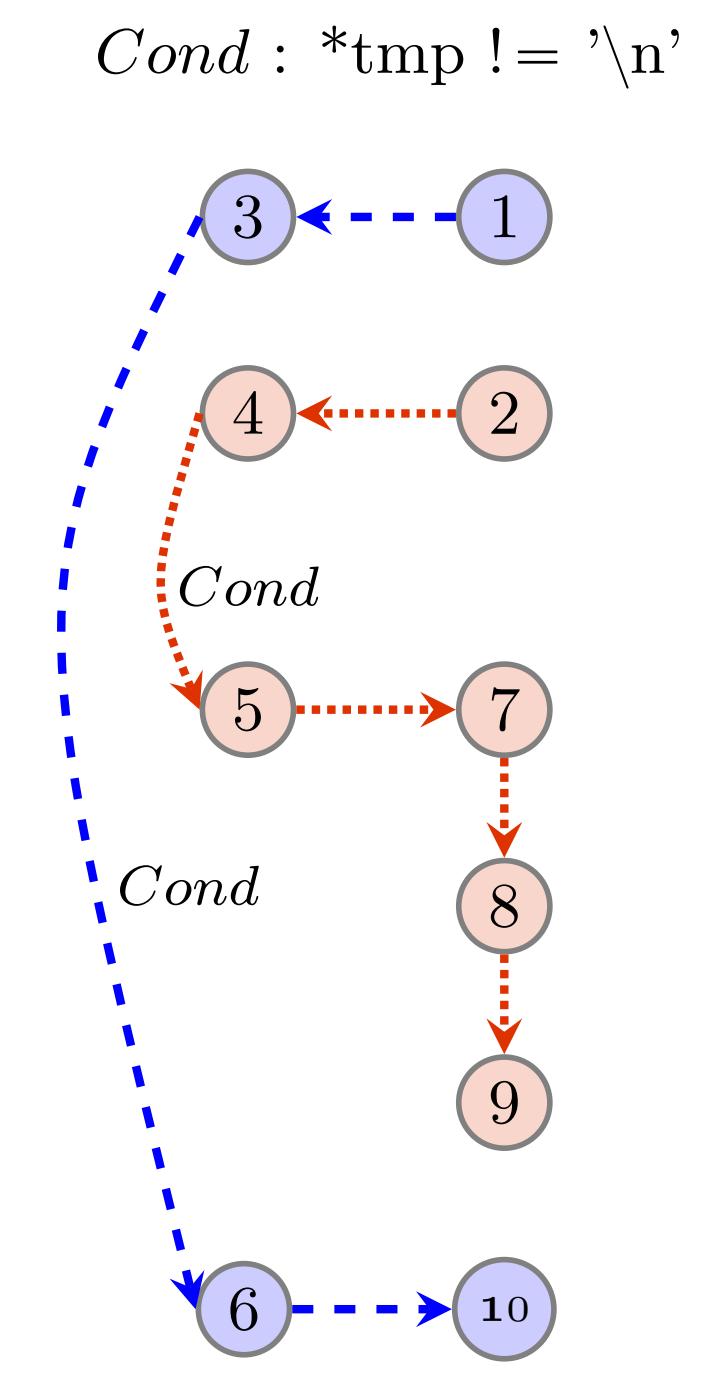
1 void SerialReadBuf() {
2     for (n=0; n<100; n++) {
3         char** buf = readBuf();
4         char* tmp = *buf;
5         if (*tmp != '\n')
6             printf("%s", *tmp);
7         else
8             continue;
9         freeBuf(buf);
10    }
11 }
12 char** readBuf() {
13     char** mBuf = malloc(); //o
14     *mBuf = malloc(); //o'
15     //... (write into **mBuf);
16     return mBuf;
17 }
18 void freeBuf(char** fBuf) {
19     char* z = *fBuf;
20     free(z);
21     free(fBuf);
22 }

```

(a) Input program



(b) Full-sparse SSA on Inter-procedural CFG



(c) SVFG (with its unlabelled edges being guarded by true)

This example, which is adapted from a real scenario in wine, demonstrates full-sparse value-flow analysis for leak detection. In this program, **readBuf** is called in a for loop in **SerialReadBuf**. Every time when **readBuf** is called, a single-char buffer formed by two objects is created: *o* at line 13 and *o'* at line 14. There are two cases. If the buffer contains a char that is not '\n', the char is printed and then both *o* and *o'* are freed. Otherwise, both objects are leaked.

## EXPERIMENTAL RESULTS

### Real Bugs Found By Saber

Code snippets from *wine-0.9.24* and *icecast-2.3.1* illustrating relevant leaky code. Arrows point from the code to specific memory operations (malloc, free, etc.) where leaks are detected.

Two scenarios with conditional leaks requiring value-flow analysis for address-taken variables: wine-0.9.24 is a tool that allows windows applications to run on Linux and icecast-2.3.1 is a streaming media server.

